### **Summit Burst Buffer**

*Chris Zimmer*







- Spider-3 center-wide GPFS parallel file system (PFS)
	- $-$  250 PB @ 2.5 TB/s
	- $-$  ~540 MB/s write performance per node when all nodes are writing

- 4,608 nodes with NVMe SSDs (Samsung PM1725a)
	- 7.3 PB Total
	- 9.67 TB/s aggregate write 27 TB/s aggregate read performance when using all nodes





• Meant to absorb bursts of I/O traffic

- What it is on Summit
	- A high performance storage device 800 GB allocated for jobs
	- Software infrastructure to support different uses (transparent drain)
	- LSF Integration to support data draining after a job has ended



# **BBAPI for Transparent Drain**

#### • Data will be drained from the Burst Buffer to Spider 3

- Uses NVMeF Offloading
	- Allows the device blocks to be read over the network without using the CPU on the compute node.

#### • Application writes data to the burst buffer

- Schedules the drain
- Gets back to work
- Data moves off on its own



## **BSCFS for Shared File Checkpoint**

#### • IBM is providing BSCFS

- Log structured file system using GPFS metadata
- Enables scientific application to write a single shared non-overlapping file using the node local burst buffers as cache.
- Requires directive calls in application to signal intent



## **Users can choose from multiple I/O options**

#### 1. Write to the PFS

- 2. Write to SSD, bulk copy at the end of the job
- 3. Write to SSD, async copy during the job (Spectral (BBAPI))
- 4. Use SCR to manage checkpoints, persist some, but not all
- 5. Use BSCFS (Single shared files to node-local burst buffers)



# **Things to know**

#### • The PFS is a shared resource

- Performance is based on shared access
	- If using the whole machine, it is 2.5 TBs/4600 nodes or  $\sim$  540 MB/s/node
	- If using 20% of the machine, it is 2.7 GB/s/node *if no other jobs are writing*
- Performance also depends on *ideal* write patterns (i.e. large, streaming writes)
	- GPFS' large random performance is much better (~90% of sequential) than Lustre's
- The node-local SSDs are not
	- Performance scales linearly with node count
	- Should exhibit low variability regardless of write pattern





- Level of effort required none; no code modification required
- Supports file-per-process, single shared file, or anything in between
- Lowest performance
- Highly variable performance



# **Write to SSD, bulk copy at the end of the job**

• Level of effort required – minimal; no code modification required

- Just change output directory
- Add file copy command to job script
- Only supports file-per-process or file-per-node
- Better performance
- If copy is at the end of the job, the user may be charged for the time to "drain" to the PFS while the next job starts and runs



# **Write to SSD, async copy during the job**

### • Eg. LibSpectral

- Level of effort required small; no code modification required, but…
- Need to set environment variable, link against library, and write to special directory
- At file close(), asynchronously copies file to the PFS directory specified in the environment variable
- Only supports file-per-process or file-per-node
- Better performance, less "drain" time (less time charges)



# **Use SCR to manage checkpoints, persist some, but not all**

#### • Level of effort required – moderate

- Requires code modification and linking against the library
- User chooses policy (persist all, persist 1 out of 10, etc.)
- Only supports file-per-process
- Best performance, especially when persisting less than all





• Level of required – Moderate; Code modification is required

- Hybrid solution
- Write to separate mount point
- Uses FUSE to redirect writes to local SSD, generate log file
- File close() triggers copy to PFS
- Can only read local data until moved to PFS, then reads come from PFS
- Designed to support single shared file
- Better single shared file performance than writing direct to PFS







