Standard Acts of Liaisons: Scientific Computing Success Stories



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Outline

- About Liaisons
- Liaison Role
- Success Stories



About Liaisons

- PhD-level scientists with expertise in computation
 - Astrophysics, biology, chemistry, climate, computer science, engineering, materials science, mathematics, nuclear physics, plasma physics, etc.
 - Experienced computational scientists with one thing in common
- Liaisons matched with INCITE projects based on science, mathematical, and algorithmic expertise
 - Can't always match for science first, e.g., I am not a chemist, but I am familiar with their math and algorithms
- Our motto: Whatever it takes!



Liaison Role

- Liaisons are collaborators whose unique expertise with leadership-level computers will enhance your experience and help you get more science done
- Levels of liaison support
 - Level 1: User support +
 - Level 2: Paratrooper fix a specific problem in your code, O(1 month)
 - Level 3: Embedded member of code development team and science collaborator



Liaison Role

- Typical liaison activities
 - Profile code performance, providing feedback to code team
 - Code porting
 - Implement solutions to problems experienced by application scientists
 - Advocate for users regarding tools, libraries, etc.
 - Collaborate scientifically



Scientific Computing Success Stories

- Recent successes, chosen to illustrate our typical tasks
- Many more success stories where these came from
- Next year: your code?



Code Profiling

- Code profiling: similar to energy audit of your home
 - Look at code's behavior, find bottlenecks, time sinks (leaking windows, old weather stripping)
 - Make recommendations for improvements to your code (replace windows, install a new water heater)
- Tools we typically use for this (you can try it too!):
 - Vampir and VampirTrace
 - Craypat



VampirTrace/Vampir

- VampirTrace: instrument code to produce trace files
 - Compile with VT wrapper, run code and obtain trace output files
- Vampir: use to visualize trace
 - Run in interactive job of nearly same size as job that produced trace files
 - Server on interactive job serves as analysis engine to local frontend



CrayPAT

- Package for instrumenting and tracing codes on Cray systems
- Run instrumented code to obtain overview of code behavior
- Re-run with refined profiling, to trace most important subroutines

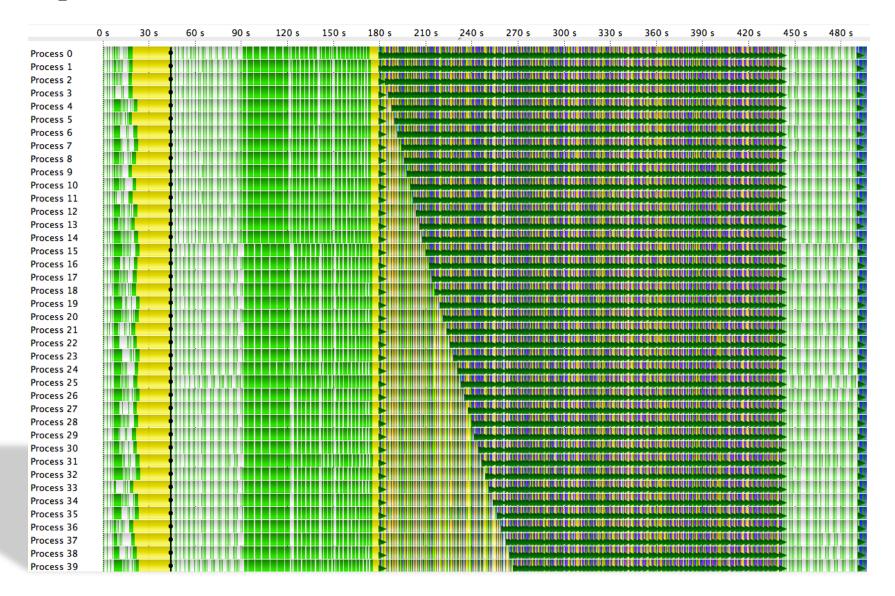


Code Profiling: Example 1 (Bigstick)

- I profiled Bigstick, a configuration interaction nuclear physics code
- Promising for large-scale nuclear configuration calculations, because it does not require as much memory
- But, does not scale well
- My task: find out why
- Used Vampir/VampirTrace to create graphical representation of code performance

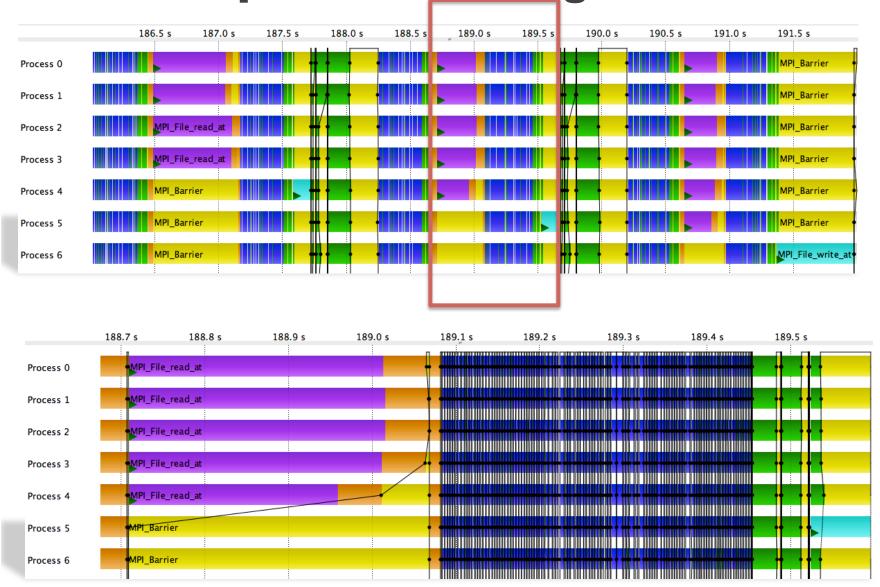


Top-Level Overview





Three Steps within Triangle





Block Reduce Phase

	213.238 s	213.240 s	213.242 s	213.244 s	213.246 s	213.248 s	213.250 s	213.252 s	213.254 s	213.256 s	213.258 s	213.260 s
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Process 33		MPI_Allreduce	MPI_Allred	duce	MPI_Allreduce	MPI_Allreduc	e MPI_	Allreduce	MPI_Allreduc	e sync	MPI_Allreduce	+ +



Bigstick: Analysis

- Triangular pattern in overview reminiscent of sequential algorithm applied across processors
 - Digging deeper shows in orthogonalization phase, processors held up by single processor writing to Lanczos vector file
 - Suggestion: reduce amount of orthogonalization performed
- Disproportionate time spent in MPI_Barrier (~30%)
 - Indicative of load imbalance
 - Barriers are within clocker subroutine, used for performance timings, obscuring evidence of load imbalance
- Majority of time in block reduce phase spent in MPI_Allreduce
 - Combining Allreduces could improve performance



Code Profiling: Example 2

- We first profiled j-coupled version of NUCCOR with CrayPat
- Discovered it spent >50% of its time sorting in test benchmark
 - Found it was using highly inefficient bubble-sort-like algorithm
 - Replaced "Frankenstein sort" with heapsort, reduced sorting to ~3% of time
- Asked collaborators what they were sorting, and why?
 - Their response: "We're sorting something?"
- Removed sorting altogether, code worked just fine, and ran 30% faster on long benchmark



Code Profiling: Example 2 (continued)

- We next profiled standard version
- Discovered it spent nearly 70% of time in single subroutine: t2_eqn_store_p_or_n
- This subroutine became focus of subsequent work with NUCCOR



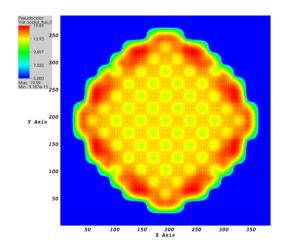
Code Porting

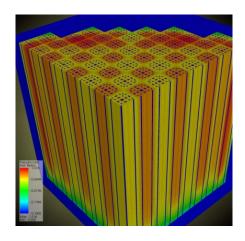
Wayne Joubert, Denovo



What is Denovo?

- Denovo is a radiation transport code used in advanced nuclear reactor design
- It solves for the density of particle flux in a 3-D spatial volume such as a reactor
- In particular, it solves the six-dimensional linear Boltzmann equation (3-space, 2-angle, 1-energy)
- Denovo scales up to 200K cores on ORNL's 2.3PF Jaguar system.
- It is part of the CASL project (Consortium for Advanced Simulation of Light Water Reactors) and the SCALE code system (Standardized Computer Analyses for Licensing Evaluation)
- It was selected as an early port code for Titan

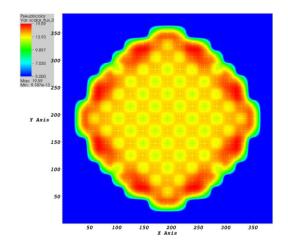


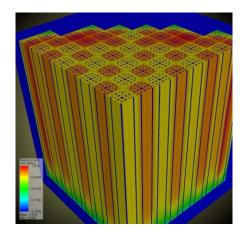




Denovo Algorithms

- Primary algorithms: the discrete ordinates method, 3-D sweep, GMRES linear solver and various eigensolvers, e.g., Arnoldi
- The execution time profile has a very prominent peak: nearly all the execution time (80-99%) is spent in a 3-D sweep algorithm.
- Because of this, the 3-D sweep must be the central focus of any effort to port Denovo to a accelerator-based system
- However, the sweep is a complex algorithm that is difficult to parallelize efficiently.

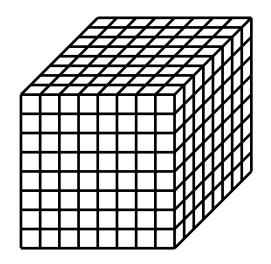


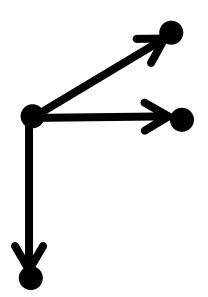




3-D Sweep Algorithm: Description

- Denovo is based on a 3-D structured grid
- The data dependency for the sweep operation is specified by a 4-point stencil
- The result at every gridcell is dependent on the result at the immediately lower gridcells in X, Y and Z.
- This induces a wavefront computation pattern a sequence of diagonal planes sweeping in from a corner.
- Thus, results at the far side of the grid cannot be computed until results at the near side are completed
- For standard parallel grid decompositions, most of the processors will be idle much of the time







How to Program the Sweep on the GPU?

- Decide what language / parallel API to use to program the GPU.
- Options:
 - 1. CUDA: a minor extension of C/C++ for GPU thread programming, also available for Fortran 90
 - 2. OpenCL: a multi-vendor standard similar to CUDA
 - 3. Compiler directives: similar to OpenMP (PGI, CAPS, Cray, ...)
- Sweep is a complex algorithm, with many dimensions. Directives may not be flexible enough or expose enough hardware functionality to get the needed performance.
- NVIDIA support OpenCL, but going forward CUDA will be better supported and more in-sync with new hardware features.
- Thus use CUDA. Use C++ for consistency with Denovo base language.



Mapping the Algorithm to the GPU

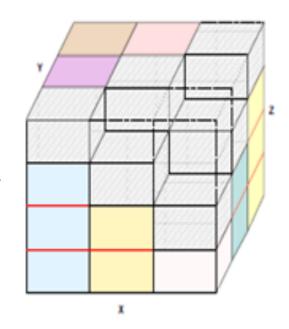
We have many candidate dimensions for parallelism: space (3), energy, moment/angle, octant, and also unknown (4 unknowns per gridcell for this discretization).

We are told by NVIDIA that we need 4K-8K threads for the GPU to keep all GPU streaming multiprocessors busy and cover various latencies.



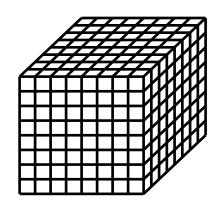
Also must have good memory access patterns (reuse of data loaded from global memory, coalesced stride-1 memory references, good use of registers, shared memory, caches on the GPU).

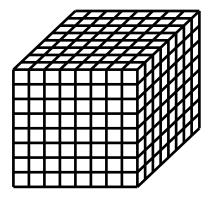
Approach: explore each problem dimension for potential thread parallelism.

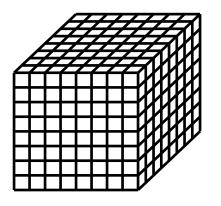


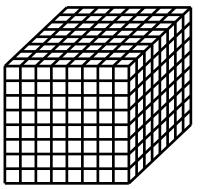
1. Parallelism in Energy

- Denovo exposes energy as a parallel dimension.
- Values for different energies are entirely independent in the 3-D sweep, thus the algorithm is embarrassingly parallel along this axis (!).
- Model problem has 256 energy groups this helps, but we need to look further in order to get to 4K-8K threads.
- Also need to use some of this 256 for node parallelism.





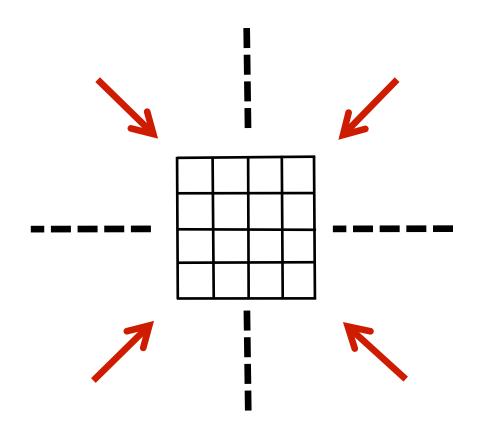






2. Parallelism in Octant

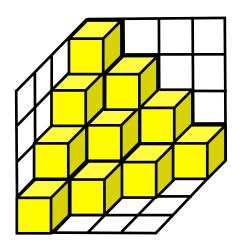
- Algorithm requires sweeps from 8 different directions.
- Sweep directions are independent, thus another 8X thread parallelism (!).
- Small issue: different octants update the same output vector, so we need to schedule properly to avoid write conflicts.





3. Parallelism in Space

- We have this recursion, as mentioned before, that makes the computations non-independent.
- However, the global KBA algorithm can be applied at this small scale (!).
- Set up block wavefronts, assign blocks to threads.
- Sync between block wavefronts.





Performance

- With this parallelization scheme, code performed at only about 1% of peak flop rate, much lower than predicted by the performance model
- Reason: excessive use of registers caused spillage to main memory, thus poor performance
- Needed to find more/better axes of parallelism



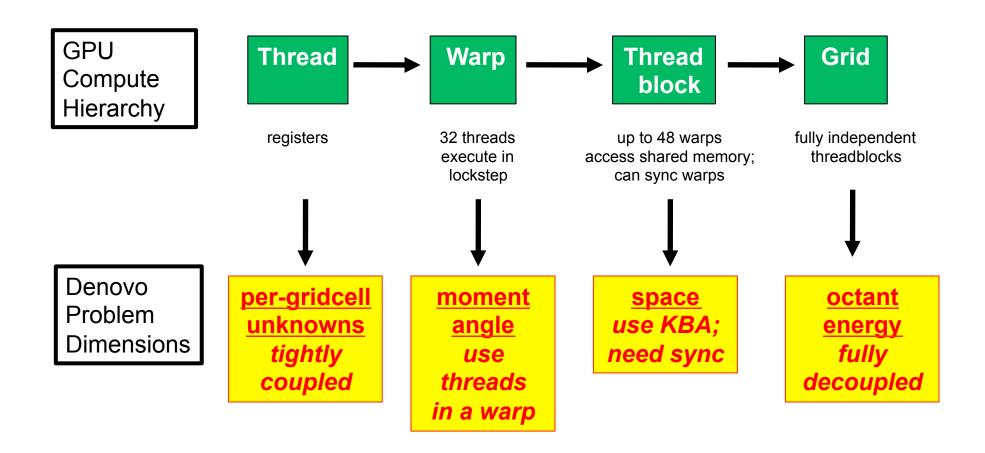
4. Parallelism in Angle, Moment

- A new strategy to parallelize the moment/angle axes at the gridcell level

 map these axes to CUDA threads in-warp.
- Small dense matrix-vector products are perfect for thread parallelism store vector in shared memory, relieve the register pressure.
- The two small matrices are the same across all gridcells (!), so they can be retained in L1 cache, to reduce a potentially high source of memory traffic.



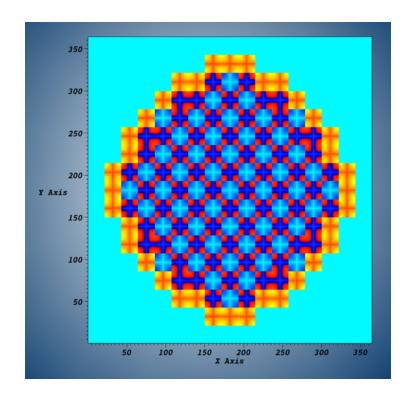
Summary of Mapping of Dimensions





Results: Test Problem

- 32x32x128 gridcells
- 16 energy groups
- 16 moments
- 256 angles
- Linear discontinuous elements 4 unknowns per gridcell



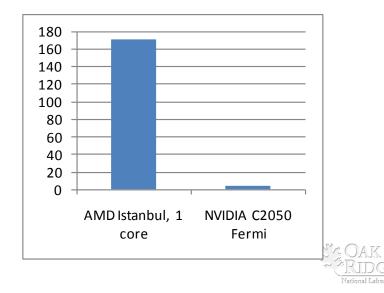


Results: Sweep GPU Performance

• Single core (AMD Istanbul) / single GPU (Fermi C2050) comparison

	AMD Istanbul 1 core	NVIDIA C2050 Fermi	Ratio
Kernel compute time	171 sec	3.2 sec	<u>54X</u>
PCIe-2 time (faces)		1.1 sec	
TOTAL	171 sec	4.2 sec	<u>40X</u>

NVIDIA Fermi is 40X faster than single Opteron core



Observations

- 40X faster than Istanbul core.
- Istanbul is 6-core, so Fermi about 7X faster than the entire Istanbul processor.
- For both CPU and GPU, code attains about 10% of peak flop rate this is considered good for this algorithm.
- Expect more optimizations to be possible going forward.



Solving Code Problems

Valentine and VPIC



VPIC

- 3-D electromagnetic relativistic particle-in-cell simulation code
- Ticket came in about VPIC failing at high core counts
- I asked newest colleague Valentine Anantharaj to take a look, with support from two experienced liaisons (Hai Ah Nam and myself)
- With Hai Ah's help, he isolated location of failure
- Libraries were incompatible because of version changes; recompiling with fresh environment and proper modules loaded eliminated the error



User Advocates

- Markus and Vasp
- Queue priority
- Software needs



VASP

- Vienna Ab-initio Simulation Package
- Licensed software, builds maintained by Markus Eisenbach
- Optimized builds that run more efficiently on Jaguar



Queue Priority

- Liaisons have science-area expertise, and understand unique needs of users
- Advocate and explain why exception is needed
- Advocate for users whose allocations have been exhausted



Software Needs

- Advocate for users' software needs
- Serve as translator between projects and Resource Utilization Council (RUC)

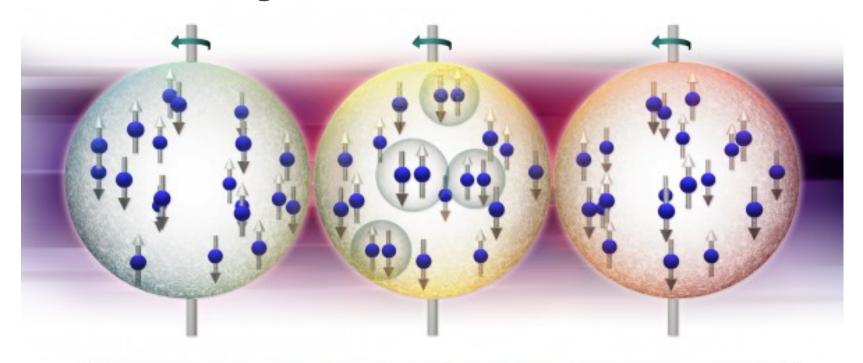


Scientific Collaboration

• Hai Ah Nam, Nuclear Physics



Nuclear Physics

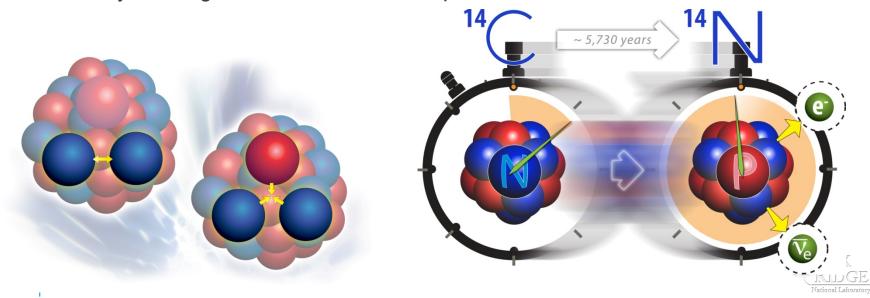


- Germanium-72 experiences at least 2 phase transitions
- Phase transitions a function of pairing of neutrons, rotation of nucleus, and temperature
- Reminiscent of superconductors



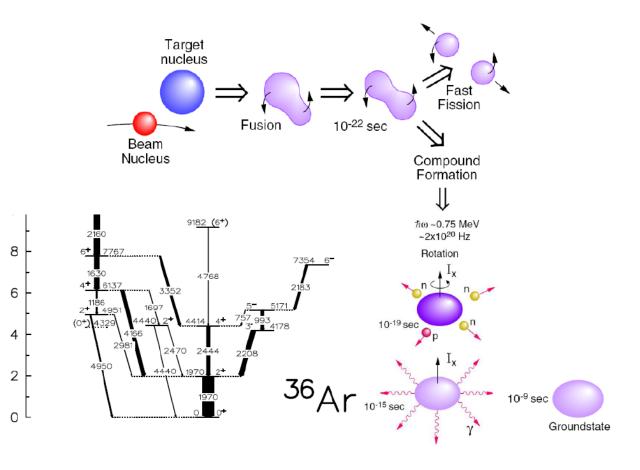
Nuclear Physics

- Studying Carbon-14 half-life with three-body interactions
 - Easy and cheap to approximate nuclear forces with two-body interactions (p+p, p+n, n+n)
 - For better accuracy, need three-body interactions (3p, 2p+1n, 1p+2n, 3n)
 using 2-body only underestimates half-life
 - BUT, adding 3-body interactions adds significant computational cost and memory requirements
 - Fortunately, on Jaguar this calculation is possible



Visualizing Spontaneous Fission of Heavy Nuclei

- Collaboration with A. Staszczak, UT/ Warsaw
- Project of summer student Elizabeth Morris
- In collaboration with visualization liaisons Dave Pugmire, Sean Ahern, Ross Toedte





Spontaneous Fission of Fermium-258

• Unusual, unexpected deformation of nucleus



Conclusions

- Liaisons are a valuable resource for INCITE projects
- Liaisons have unique HPC skills that projects can take advantage of
- Whatever it takes!



Acknowledgments

- Wayne Joubert
- Hai Ah Nam
- Scientific Computing group
- INCITE project participants



Questions?



